ECE 587 – Hardware/Software Co-Design Lecture 03 State-Based Models I

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- \blacktriangleright This lecture: 3.1, 3.1.2
- ▶ Next lecture: 3.1.2

Outline

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Models of Computation

- ▶ Any non-trivial functionality must involve some kind of computation.
- \blacktriangleright It is beneficial to specify the functionality just at the abstraction level of the computation.
	- \blacktriangleright It's intuitive.
	- ▶ Computations are behavioral. No implementation detail is necessary.
	- ▶ Computations are based on mathematics. There may exist tools to automate the remaining design process.
- ▶ Models of Computation (MoCs)
	- \triangleright Serve as basis to reason about computation/constraints
	- ▶ Utilize formal language, e.g. certain kind of mathematics
	- ▶ May have different supported features, complexity, and expressive power.
- \triangleright MoCs define computations by specifying when to perform operations.
	- \blacktriangleright The time here is not absolute time but relative ordering.
	- ▶ So ultimately it depends on how synchronizations are employed.
- ▶ Fully synchronized model: Finite State Machine
- ▶ Fully ordered without synchronization: Sequential Programs
- ▶ No synchronization at all: Dataflow
- ▶ We will first focus on FSM and move to other models in the next few weeks.

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Finite-State Machine (FSM)

$< S, I, O, f, h >$

- \blacktriangleright Set of states S
- \triangleright Set of input symbols I
- \triangleright Set of output symbols O
- ▶ Next-state function $f : S \times I \rightarrow S$
- ▶ Output function $h: S \times I \rightarrow O$
- \triangleright Some systems may specify initial states and/or final states
- ▶ Encoding of states and input/output symbols in HW/SW
	- ▶ This condition will sometimes be relaxed so one can handle extremely large systems.
- \blacktriangleright Implementation of f and h in HW/SW

Representations of FSM

\blacktriangleright Graph representation

- ▶ States as vertices
- \triangleright State transitions as edges (annotated with inputs/outputs)
- \blacktriangleright Intuitive, but if there are too many possible states, it becomes unmanageable.
- ▶ Functional representation
	- \blacktriangleright If one can efficiently specify f and h, then the FSM can be simulated from any initial state and a trace of inputs, fulfilling most computational tasks.
	- \triangleright Can handle extremely large systems
- ▶ Since a FSM has a finite number of possible states, one can represent, or encode, a state using a fixed number of bits.
	- \blacktriangleright E.g. if there are 16 possible states, a 4-bit encoding can be applied.
- ▶ Similarly you can encode inputs and outputs.
- \triangleright Under such encodings, the functions f and h become boolean functions.

FSM vs. Register Transfer Level (RTL)

- \blacktriangleright That's exactly how RTL is defined.
	- \blacktriangleright Just change the state bits to registers
- \blacktriangleright The key here is encoding.
	- ▶ Encoding enables us to specify extremely large FSMs.
	- ▶ Different encodings may lead to specifications with different complexity, though for system design we prefer to use the most intuitive one.
- ▶ We will still distinguish functional representations of FSM from RTL as they have different purposes.
	- \blacktriangleright Though mathematically there is no difference.

Implement FSMs

▶ Hardware: as Synchronous Circuits

- ▶ Utilize the connection between functional representation and RTL
- \blacktriangleright Exactly one state transition happens per clock cycle.
- \blacktriangleright High speed/low power/energy consumption
- ▶ Usually known as cycle-accurate behavior
- ▶ Software: follow either graph or functional representations
	- ▶ Tedious, better to have tools to generate code
	- ▶ Not efficient in both time and power
	- ▶ But is a very powerful architecture to build complex software that needs to react to external events, e.g. networking and graphical user interface.

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Input Validation

- ▶ Consider an application that requires to validate user inputed numbers
	- \triangleright Assume the input is a character string
	- \blacktriangleright End of string must be enter.
- \blacktriangleright A valid integer
	- \blacktriangleright If the first character is not a digit, then it must be either $+$ or −.
	- \triangleright Except the first character and the ending enter, all characters are digits.
	- \blacktriangleright The most significant digit must not be 0.
	- ▶ The integer may contain arbitrary number of digits.
- ▶ Additional tasks
	- \blacktriangleright Deal with floating-point numbers
	- \blacktriangleright Extract the number during validation
	- \blacktriangleright Implement the designs in a programming language.
- \blacktriangleright How to approach this or similar problems?

A Simple FSM

- ▶ How does it work?
	- ▶ Starting from S0
	- ▶ Process exactly one character per transition.
- \blacktriangleright This simple example accepts numbers like 1.2, 4.5, but not 11 or 1.21.

A More Complex FSM

- ▶ Build a FSM to recognize integers.
- ▶ Extend it to handle floating-point numbers.

Extract Numbers

▶ Focus on integers but make it easy to extend our solution to floating-point numbers etc.

Software Implementation

```
enum {S0, S1, ..., OK, FAIL};
int state = S0, sign = 1, num = 0;
while ((S0 != OK) & kk (SO != FAIL)) {
  int next_state = state; // assume state remain the same by default
  int ch = read_one_input();
  if (\text{state} == 50) {
    if (ch == '-'') {
      next\_state = S1; sign = -1;} else if (isdigit(ch)) {
      next state = S1; num = ch-'0';
    } ...
  } else if (state == S1) {
 } ...
  state = next_state;
}
num *= sign;
```
- \triangleright Make use of a single loop to drive the state transitions.
- \triangleright Use two levels of branches to handle combinations of current state and input.
- \blacktriangleright It can handle any FSM no matter how complicated it is.

Discussions

- ▶ From the FSM model, it will be much easier for the designers to utilize tools at hand to implement the validation as either hardware or software.
- ▶ Such problems are special cases of Regular Expressions.
	- \blacktriangleright It is used almost everywhere when text is processed.
	- ▶ Many places require to run it very efficiently, e.g. to filter certain information from the network at realtime.
- ▶ Regular expressions can be modeled by a special kind of FSMs called nondeterministic FSM.
	- ▶ There is a mapping from graph representation of nondeterministic FSM to RTL, which enable one to implement it quite efficiently in hardware.
	- ▶ The challenge in hardware implementation is reconfigurability without much overhead.
	- ▶ Software implementations are based on the same idea but are much more awkward.