

Binary Trees and Prefix Codes

Imagine that you have an alphabet of symbols

■ e.g. A, B, ..., Z, a, b, ..., z, `', `.'

- We wish to represent a string of these symbols as a string of bits
 - e.g. "This is a string of characters" becomes "01100111000101010111001"

Method 1: Use a fixed number for each symbol

◆ Map A -> 1, B -> 2, ...

- This is a string of characters" becomes a string of numbers
 - "20,34,35,..."
- I required commas to separate the characters!
 - Use a fixed width, padded with leading 0's instead
 - "020034035..."

Method 1: Fixed width

- How wide to my characters need to be
 - Using a string of bits
 - How many bits are needed to represent the largest character?
- \bullet ceil(log₂(n)) bits
- Use that many bits for each character
- 4
- This is the system used within the computer with 8 bits for each ASCII code

Method II: Prefix codes

- For each symbol, we'll use a code with a special property
 - No code is the prefix of any other code
- How does this work?
- Decoding:
 - We read in the codes one bit at a time
 - When we have a code we recognise, it must be the end of a symbol
 - It cannot be part of a longer symbol because no code is the prefix of another code



Making a Prefix Code

- We want the code to be efficient
 - No strings longer than necessary
 - No wasted strings
- A code is a set of strings of binary digits, such that no string corresponding to one symbol is the prefix of a string corresponding to another symbol
- ◆ In a tree, leaf nodes have no children
 - No path from the root to a leaf is the prefix of a path from the root to another node

Binary Trees and Prefix Codes

Binary trees are in one to one correspondence with Prefix Codes

A:00, B:010, C:011, D:10, E:11



Prefix Trees

- Binary Trees
 - The left child corresponds to 0, the right to



- Each leaf contains a symbol
- The code for a symbol corresponds to the path from the root to the leaf containing that symbol

Encoding and Decoding

Imagine the encoder and decoder running in parallel

Encoding

- Start from the root
- While you are not at the symbol's leaf
 - If the symbol you wish to send is a left decendant, send 0 and move to your left child, else send 1 and move to your right child
- Decoding
 - Start from the root
 - While you are not at a leaf
 - Read a bit. If it is 0 then move to your left chile, else move to your right child



































Back to Method I: Balanced Tree

- Method I was to used fixed length code words
- Each path from the root to a leaf is the same length: a balanced tree
- Balanced trees are good for worst case path length. Are they good for coding?
 - Yes, if you assume the worst case
 - But we can normally do better...

Statically optimal codes

- Want common symbols to have short codes
- This will make uncommon symbols have longer codes
 - In a tree with a fixed number of leave/symbols, moving one leaf/symbol closer to the root will move others further away

Huffman codes

From Shannon's information theory, The optimal static code assigns $-\log_2(p)$ bits to a symbol that occurs with probability p



It is possible to make a Huffman code tree with this property

Will look at this later in the course

Adaptive Codes

As long as the same change is made in both sending and receiving trees/codes, there is no reason why the tree/code must remain static

- Send a character using the initial tree
- Update the tree using that character
 - Can also be updated in the receiver as it already has the character
- Send the next character















