## Trees and Codes



## Binary Trees and Prefix Codes

- Imagine that you have an alphabet of symbols
- e.g. A, B, ..., Z, a, b, ... z, ' ', '.'
$\bullet$ We wish to represent a string of these symbols as a string of bits
- e.g. "This is a string of characters" becomes "01100111000101010111001"


## Method 1: Use a fixed number for each symbol

* Map A -> 1, B -> 2, ...
* "This is a string of characters" becomes a string of numbers
- "20,34,35,..."
- I required commas to separate the characters!
- Use a fixed width, padded with leading 0's instead - "020034035..."


## Method 1: Fixed width

- How wide to my characters need to be
- Using a string of bits
- How many bits are needed to represent the largest character?
- ceil( $\left.\log _{2}(n)\right)$ bits
- Use that many bits for each character
- This is the system used within the computer with 8 bits for each ASCII code


## Method II: Prefix codes

- For each symbol, we'll use a code with a special property
- No code is the prefix of any other code

How does this work?

- Decoding:
- We read in the codes one bit at a time
- When we have a code we recognise, it must be the end of a symbol
- It cannot be part of a longer symbol because no code is the prefix of another code


## Method II: Prefix codes

- Example
- A:00, B:010, C:011, D:10, E:11
- ADBECABADE
*00100101101100010001011
- A D B E C A B A D E
- The string decodes


## Making a Prefix Code

- We want the code to be efficient
- No strings longer than necessary
- No wasted strings
- A code is a set of strings of binary digits, such that no string corresponding to one symbol is the prefix of a string corresponding to another symbol
- In a tree, leaf nodes have no children
- No path from the root to a leaf is the prefix of a path from the root to another node


## Binary Trees and Prefix Codes

- Binary trees are in one to one correspondence with Prefix Codes
- A:00, B:010, C:011, D:10, E:11



## Prefix Trees

- Binary Trees
- The left child corresponds to 0 , the right to 1
- Each leaf contains a symbol
*The code for a symbol corresponds to the path from the root to the leaf containing that symbol


## Encoding and Decoding

- Imagine the encoder and decoder running in parallel
- Encoding
- Start from the root
- While you are not at the symbol's leaf
- If the symbol you wish to send is a left decendant, send 0 and move to your left child, else send 1 and move to your right child
- Decoding
- Start from the root
- While you are not at a leaf
- Read a bit. If it is 0 then move to your left chile, else move to your right child


## Encoding and Decoding: ACD

- Encoding:ACD
- Decoding:



## Encoding and Decoding: ACD

- Encoding:ACD
- Decoding:



## Encoding and Decoding: ACD

- Encoding:ACD
- Decoding:



## Encoding and Decoding: ACD

- Encoding:ACD
- Decoding:



## Encoding and Decoding: ACD

- Encoding:CD
- Decoding:A



## Encoding and Decoding: ACD

- Encoding:CD
- Decoding:A



## Encoding and Decoding: ACD

- Encoding:CD
- Decoding:A



## Encoding and Decoding: ACD

- Encoding:CD
- Decoding:A



## Encoding and Decoding: ACD

- Encoding:CD
- Decoding:A



## Encoding and Decoding: ACD

- Encoding:CD
- Decoding:A



## Encoding and Decoding: ACD

- Encoding:CD
- Decoding:A



## Encoding and Decoding: ACD

- Encoding:D
- Decoding:AC



## Encoding and Decoding: ACD

- Encoding:D
- Decoding:AC



## Encoding and Decoding: ACD

- Encoding:D
- Decoding:AC



## Encoding and Decoding: ACD

- Encoding:D
- Decoding:AC



## Encoding and Decoding: ACD

- Encoding:D
- Decoding:AC



## Encoding and Decoding: ACD

* Encoding:
* Decoding:ACD



## Back to Method I: Balanced Tree

- Method I was to used fixed length code words
- Each path from the root to a leaf is the same length: a balanced tree
- Balanced trees are good for worst case path length. Are they good for coding?
- Yes, if you assume the worst case
- But we can normally do better...


## Statically optimal codes

- Want common symbols to have short codes
*This will make uncommon symbols have longer codes
- In a tree with a fixed number of leave/symbols, moving one leaf/symbol closer to the root will move others further away


## Huffman codes

*From Shannon's information theory, The optimal static code assigns $-\log _{2}(p)$ bits to a symbol that occurs with probability $p$
$\bullet$ It is possible to make a Huffman code tree with this property

- Will look at this later in the course


## Adaptive Codes

*As long as the same change is made in both sending and receiving trees/codes, there is no reason why the tree/code must remain static

- Send a character using the initial tree
- Update the tree using that character
- Can also be updated in the receiver as it already has the character
- Send the next character


## Encoding and Decoding: ACD

- Encoding:ACD
- Decoding:



## Encoding and Decoding: ACD

- Encoding:CD
- Decoding:A



## Encoding and Decoding: ACD

- Encoding:CD
- Decoding:A

Make same change in both trees:
Rotate A's parent


## Encoding and Decoding: ACD

- Encoding:CD
- Decoding:A



## Encoding and Decoding: ACD

- Encoding:D
- Decoding:CA



## Encoding and Decoding: ACD

- Encoding:D
- Decoding:CA



## Encoding and Decoding: ACD

- Encoding:D
- Decoding:CA



## Encoding and Decoding: ACD

* Encoding:
$\bullet$ Decoding:CAD


